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A New Differential Formalism for Interval-Valued Functions
and Its Potential Use in Detecting 1-D Landscape Features

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Abstract: In many practical problems, it is important to know the slope (derivative) \( \frac{dy}{dx} \) of one quantity \( y \) with respect to some other quantity \( x \). For example, different 1-D landscape features can be characterized by different values of the derivative \( \frac{dy}{dx} \), where \( y \) is an altitude, and \( x \) is a horizontal coordinate. In practice, we often know the values of \( y(x) \) for different \( x \) with interval uncertainty. How can we then find the set of possible values of the slope? In this paper, we formulate this problem of differentiating interval-values functions in precise terms, and we describe an (asymptotically) optimal algorithm for computing the corresponding derivative.

Keywords: interval computations, differentiation, landscape

1 Introduction
In many real-life problems, we want to know the values of the derivatives. In many areas of science and engineering, we are interested in slopes. For example, a 1-D landscape is described as a dependence of the altitude \( y \) on the coordinate \( x \); different 1-D landscape features are defined by different values of the slope \( \frac{dy}{dx} \) of this dependence: low values of this slope correspond to a plain, high values to steep mountains, and medium values to a hilly terrain.

Interval uncertainty. In the ideal situation, when we know the exact values of \( y(x) \) for every \( x \), we can simply differentiate the corresponding dependence. In practice, however, the information on \( y \) comes from measurements, and measurements are never exact. E.g., in the landscape example, we measure the altitudes \( y_1, \ldots, y_n \) at different points \( x_1 < \ldots < x_n \). Since the measurements are not exact, the measured values \( \bar{y}_i \) are, in general, slightly different from the the (unknown) actual altitudes \( y_i \).

For measuring instruments, we usually have an upper bound \( \Delta_i \) on the measurement error

\[
\Delta y_i \stackrel{\text{def}}{=} \bar{y}_i - y_i.
\]

This upper bound is usually provided by the manufacturer of this instrument: \( |\Delta y_i| \leq \Delta_i \). Thus, after the measurement, the only information that we have about the actual (unknown) value \( y_i \) is that this value belongs to the interval \( y_i = [\underline{y}_i, \bar{y}_i] \), where \( \underline{y}_i \stackrel{\text{def}}{=} \bar{y}_i - \Delta_i \) and \( \bar{y}_i \stackrel{\text{def}}{=} \bar{y}_i + \Delta_i \) (for a more detailed description of interval uncertainty, see, e.g., [6, 7, 8, 11]).

Thus, the only information that we have about the actual dependence \( y = f(x) \) of \( y \) on \( x \) is that the (unknown) function \( f(x) \) belongs to the class

\[
F \stackrel{\text{def}}{=} \{ f(x) \mid f(x_i) \in y_i \text{ for all } i = 1, \ldots, n \}.
\] (1)

We also know that the (unknown) function \( f(x) \) is smooth (differentiable) – because otherwise, the notion of a slope does not make sense.

In many practical applications, the derivative has a physical meaning, and this meaning implies that it is itself a continuous (or even differentiable) function. For example, when we monitor the locations \( y_i \) of a particle at different moments of time \( x_i \), then the
very fast when no one was looking (i.e., in between
In other words, for every interval
always possible to have a function
known) function
look for areas where the slope takes a given value
somewhere, the slope was indeed equal to
What we are really interested in is not whether it is
imply
Comment. The notation $F'(\{a,b\})$ looks like the no-
tation of a range for a real-valued function, but it is
not a range: in contrast to range, if an interval is nar-
row enough, we can have $F'(\{a,b\}) = \emptyset$ (see exam-
pies below).
This newly defined derivative does share some prop-
erties of the range. For example, it is well known
that the range is inclusion-monotonic – in the sense
that $[a,b] \subseteq [c,d]$ implies $f'([a,b]) \subseteq f'([c,d])$.
From this property of the range, we can conclude that
$a,b \subseteq [c,d]$ implies $F'([a,b]) \subseteq F'([c,d])$ – i.e.,
that the newly defined derivative is also inclusion-
monotonic. Thus, if the union $A \cup B$ of two in-
tervals is also an interval, we have $F'(A \cup B) \supseteq
F'(A) \cup F'(B)$.

Formulation of the problem. How can we compute
the derivative of an interval function? The above def-
ition, if taken literally, requires that we consider all
(infinitely many) functions $f \in F$ – which is compu-
tationally excessive. Thus, we must find an efficient
algorithm for computing this derivative. This is what
we will do in this paper.

We will try our best to make sure that these algo-
rimds are not simply tricks, that the ideas behind
these algorithms are clear and understandable.
Therefore, instead of simply presenting the final al-
gorithm, we will, instead, present our reasoning in a
series of auxiliary results that eventually leads to the

2 Precise Formulation of the Problem

Definition 1. By an interval function $F$, we mean
a finite sequence of pairs $\langle x_i, y_i \rangle$ ($i = 1, 2, \ldots, n$),
where for each $i$, $x_i$ is a real number, $y_i$ is a non-
degenerate interval, and $x_1 < x_2 < \ldots < x_n$.

Definition 2. We say that a function $f : R \rightarrow R$
from reals to reals belongs to an interval function
$F = \{ \langle x_1, y_1 \rangle, \ldots, \langle x_n, y_n \rangle \}$ if $f(x)$ is continu-
ously differentiable and for every $i$ from 1 to $n$, we have $f(x_i) \in y_i$.

Definition 3. Let $F$ be an interval function, and let
$[a, b]$ be an interval. By a derivative $F'(\{a,b\})$, we mean the intersection

$$F'(\{a,b\}) \overset{\text{def}}{=} \bigcap_{f \in F} f'(\{a,b\}),$$

where $f'(x)$ denotes the derivative of a differentiable
function $f(x)$, and $f'(\{a,b\}) \overset{\text{def}}{=} \{ f'(x) \mid x \in \{a,b\} \}$
is the range of the derivative $f'(x)$ over the interval
$[a,b]$.

In other words, we say that the slope is guaranteed to
attain a given value $s$ somewhere on a given interval
$[a,b]$ if for every function $f \in F$, the range $f'([a,b])$
of its derivative $f'(x)$ contains the value $s$. In mathe-
matical terms, this means that the value $s$ belongs to
the intersection of the ranges $f'([a,b])$ corresponding
to all $f \in F$.

This intersection thus describes the “range of the
derivative” of the interval function $F$ on the given
interval $[a,b]$. In other words, we arrive at the fol-
lowing definitions.
asymptotically optimal algorithms for computing the desired derivative \(F'(a, b)\).

**Previous work.** In our research, we were guided by results from two related research directions:

First, we were guided by different definitions of differentiation of an interval function that have been proposed by interval computations community [2, 7, 10, 12, 13, 15, 16, 17, 18, 19]. The main difference from our problem is that most of these papers assume that we have intervals \([a, b]\) for all \(x\), while we consider a more realistic situation when the interval bounds on \(f(x)\) are only known for finitely many values \(x_1, \ldots, x_n\).

Second, we were guided by a paper [20] in which an algorithm was developed to check for local maxima and minima of an interval function \(f\). This result has been applied to detecting geological areas [1, 3, 4] and to financial analysis [5]. This result can be viewed as detecting the areas where the derivative is equal to 0 – and, in this sense, as a particular case of our current problem.

### 3 First Auxiliary Result: Checking Monotonicity

**Definition 4.** We say that a function \(f(x)\) is strongly increasing if \(f'(x) > 0\) for all \(x\).

**Comment.** Every strongly increasing function is strictly increasing, but the inverse is not necessarily true: the function \(f(x) = x^3\) is strictly increasing but not strongly increasing.

**Proposition 1.** For every interval function \(F\), the existence of a strongly increasing function \(f \in F\) with \(f'(x) > 0\) is equivalent to

\[
y_i < \overline{y}_j \quad \text{for all } i < j.
\]

**Proof.** If \(f \in F\) and \(f(x)\) is strongly increasing, then it is also strictly increasing hence for every \(i < j\), the inequality \(x_i < x_j\) implies that \(f(x_i) < f(x_j)\). Since \(f \in F\), we have \(f(x_i) \in y_i = [\underline{y}_i, \overline{y}_i]\) and \(f(x_j) \in y_j = [\underline{y}_j, \overline{y}_j]\). Thus, \(\underline{y}_i \leq f(x_i) < f(x_j) \leq \overline{y}_j\), we conclude that \(\underline{y}_i < \overline{y}_j\), which is exactly the inequality (2).

Vice versa, let us assume that the inequalities (2) are satisfied, and let us design the corresponding strictly increasing function \(f \in F\). We will first build a piece-wise linear strictly increasing function \(f_0(x)\) for which \(f_0(x_i) \in y_i\), and then we will show how to modify \(f_0(x)\) into a continuously differentiable strongly increasing function \(f \in F\).

According to the inequalities (2), all the differences \(\overline{y}_j - \underline{y}_i\) \((i < j)\) are positive. Since all intervals are non-degenerate, the differences \(\overline{y}_i - \underline{y}_i\) are also positive. Let us denote the smallest of these positive numbers \(\Delta\). For every \(i\), let us denote

\[
y_i \overset{\text{def}}{=} \max(\underline{y}_1, \ldots, \underline{y}_i) + \frac{i}{2n} \cdot \Delta. \tag{3}
\]

We will then design \(f_0(x)\) as a piece-wise linear function for which \(f_0(x_i) = y_i\). To show that \(f_0(x)\) is the desired piece-wise linear function, we must show that for every \(i, y_i \in y_i\), and that this function is strictly increasing, i.e., that \(i < j\) implies \(y_i < y_j\).

That \(i < j\) implies \(y_i < y_j\) is clear: the first (maximum) term in the formula (3) can only increase (or stay the same) when we replace \(i\) by \(j\), and the second term increases. Thus, it is sufficient to prove that \(y_i \in y_i = [\underline{y}_i, \overline{y}_i]\), i.e., that \(\underline{y}_i \leq y_i\) and \(y_i \leq \overline{y}_i\). We will actually prove a stronger statement: that \(\overline{y}_i < y_i\) and \(y_i < \overline{y}_i\).

The first inequality \(y_i < \overline{y}_i\) follows directly from the formula (3): by definition of a maximum, \(\max(\underline{y}_1, \ldots, \underline{y}_i) \geq \underline{y}_i\), and when we add a positive number to this maximum, the result only increases. So, \(y_i\) is actually larger than \(\underline{y}_i\).

Let us now prove that \(y_i < \overline{y}_i\). Indeed, by definition of \(\Delta\), for all \(k \leq i\), we have \(\underline{y}_k + \Delta \leq \overline{y}_i\), hence (since \((i/2n) \cdot \Delta < \Delta\)) \(\underline{y}_i + (i/2n) \Delta < \overline{y}_i\). Thus, \(y_i\) – which is the largest of the values \(\overline{y}_i\) – is also smaller than \(\overline{y}_i\). So, the desired \(f_0(x)\) is designed.

Let us now show how to build the corresponding continuously differentiable function \(f(x)\). For the piece-wise linear function \(f_0(x)\), the first derivative \(f_0'(x)\) is piece-wise constant; since the function \(f_0(x)\) is strictly increasing, the values \(f_0'(x)\) are all positive. Around each discontinuity point \(x_i\), replace the abrupt transition with a linear one; as we integrate the resulting function, we get a new function \(f(x)\) that is continuously differentiable and – since the new values of the derivative are still everywhere positive – strongly increasing. When the replacement is fast enough, the change in the value \(f'(x)\) is so small that \(f'(x_i)\) is still inside the desired interval \(y_i\). The proposition is proven.

Similarly, we can prove the following results:
Definition 5. We say that a function $f(x)$ is strongly decreasing if $f'(x) < 0$ for all $x$.

Proposition 2. For every interval function $F$, the existence of a strongly decreasing function $f \in F$ is equivalent to

$$y_i > y_j \quad \text{for all } i < j.$$  

(4)

Proposition 3. For every interval function $F$ and for every interval $[a, b]$, the existence of a function $f \in F$ that is strongly increasing on the interval $[a, b]$ is equivalent to

$$y_i < y_j \quad \text{for all } i < j$$

for which $x_i, x_j \in [a, b]$.  

(5)

Proposition 4. For every interval function $F$ and for every interval $[a, b]$, the existence of a function $f \in F$ that is strongly decreasing on the interval $[a, b]$ is equivalent to

$$y_i > y_j \quad \text{for all } i < j$$

for which $x_i, x_j \in [a, b]$.  

(6)

4 Second Auxiliary Result: Checking Whether $0 \in F'([a, b])$

Proposition 5. For every interval function $F$ and for every interval $[a, b]$, $0 \in F'([a, b])$ if and only if neither conditions (5) nor conditions (6) are satisfied.

Proof. Let us first show that if either the conditions (5) or the conditions (6) are satisfied, then $0 \notin F'([a, b])$.

Indeed, according to Proposition 3, if the conditions (5) are satisfied, then there exists a function $f \in F$ that is strongly increasing on $[a, b]$. For this function, $f'(x) > 0$ for all $x \in [a, b]$; therefore, $f'([a, b]) \subseteq (0, \infty)$. Since $F'([a, b])$ is defined as the intersection of such range sets, we have $F'([a, b]) \subseteq f'([a, b]) \subseteq (0, \infty)$ hence $0 \notin F'([a, b])$.

Similarly, if the conditions (6) are not satisfied, then $0 \notin F'([a, b])$.

Vice versa, let us assume that neither the conditions (5) nor the conditions (6) are satisfied, and let us show that then $0 \notin F'([a, b])$. Indeed, let $f \in F$ be an arbitrary function from the class $F$. Since the conditions (5) are not satisfied, the function $f(x)$ cannot be strongly increasing; therefore, there must be a point $x_1 \in [a, b]$ for which $f'(x_1) \leq 0$. Similarly, since the conditions (6) are not satisfied, the function $f(x)$ cannot be strongly decreasing; therefore, there must be a point $x_2 \in [a, b]$ for which $f'(x_2) \geq 0$.

Since the function $f(x)$ is continuously differentiable, the continuous derivative $f'(x)$ must attain the 0 value somewhere on the interval $[x_1, x_2] \subseteq [a, b]$. In other words, $0 \in f'([a, b])$ for all $f \in F$. Thus, $0$ belongs to intersection $F'([a, b])$ of all possible ranges $f'([a, b])$. The proposition is proven.

5 Third Auxiliary Result and Final Description of $F'([a, b])$

Definition 6. Let

$$F = \{\langle x_1, y_1 \rangle, \ldots, \langle x_n, y_n \rangle\}$$

be an interval function, and let $v$ be a real number. Then, we define a new interval function $F - v \cdot x$ as follows:

$$F - v \cdot x = \{\langle x_1, y_1 - v \cdot x_1 \rangle, \ldots, \langle x_n, y_n - v \cdot x_n \rangle\},$$

where, for an interval $y = [y, \bar{y}]$ and for a real number $c$, the difference $y - c$ is defined as $[y - c, \bar{y} - c]$.

It is easy to prove the following auxiliary result:

Proposition 6. For every interval function $F$ and for every interval $[a, b]$, $v \in F'([a, b])$ if and only if $0 \in (F - v \cdot x)'([a, b])$.

This result leads to the following description of the derivative $F'([a, b])$:

Proposition 7. For every interval function $F$ and for every interval $[a, b]$, let $i_0$ and $j_0$ be the first and the last index of the values $x_i$ inside $[a, b]$. Then $F'(a, b) = \{F_{i_0, j_0}^+, F_{i_0, j_0}^-\}$, where

\[
E_{i_0, j_0}^+ \overset{\text{def}}{=} \min_{i_0 \leq i < j \leq j_0} \Delta_{ij},
\]

(7)

\[
F_{i_0, j_0}^- \overset{\text{def}}{=} \max_{i_0 \leq i < j \leq j_0} \Delta_{ij},
\]

(8)

\[
\Delta_{ij} \overset{\text{def}}{=} \frac{y_j - y_i}{x_j - x_i},
\]

(9)

\[
\Delta_{ij} \overset{\text{def}}{=} \frac{\bar{y}_j - \bar{y}_i}{x_j - x_i},
\]

(10)

and $[p, q] \overset{\text{def}}{=} \{x \mid p \leq x \land x \leq q\}$ – so when $p > q$, the interval $[p, q]$ is the empty set.
Comment 1. The above expression is rather intuitively reasonable because the ratios \( \Delta_{ij} \) and \( \Sigma_{ij} \) are finite differences – natural estimates for the derivatives.

Comment 2. As a corollary of this general result, we can conclude that if the interval \([a, b]\) contains a single point \(x_i\) (or no points at all), then

\[ F'(\{a, b\}) = \emptyset. \]

Mathematically, this conclusion follows from our general result because in this case, there is no pair \(i < j\), so the minimum and the maximum are taken over an empty set. By definition, the minimum of an empty set is infinite, so \( F_{\min} = +\infty\); similarly, \( F_{\max} = -\infty\). Here, \( F_{\min} > F_{\max} \), so the interval is empty. Intuitively, however, this conclusion can be understood without invoking minima and maxima over an empty set.

Indeed, let us assume that the given interval \([a, b]\) contains only one point \(x_i\) from the original list \(x_1, \ldots, x_n\). Then, for any real number \(s\), we can take, as \( f \in F\), a function that takes an arbitrary value \(y_i \in y_i\) for \(x = x_i\) and that is linear with a slope \(s\) on \([a, b]\) – i.e., the function

\[ f(x) = y_i + s \cdot (x - x_i). \]

For this function \(f(x)\), the range \(f'([a, b])\) of the derivative \(f'(x)\) on the interval \([a, b]\) consists of a single point \(s\). Thus, if we take two such functions corresponding to two different values of \(s\), then the intersection of their ranges is empty. Therefore, the range \(f'([a, b])\) – which is defined (in Definition 3) as the intersection of all such ranges \(f'([a, b])\) – is also empty.

**Proof of Proposition 7.** The fact that conditions (5) are not satisfied means that there exist value \(i_0 \leq i < j \leq j_0\) for which \(\bar{y}_i \geq \bar{y}_j\). The fact that the conditions (6) are not satisfied means that there exist values \(i_0 \leq i' < j' \leq j_0\) for which \(\bar{y}_{i'} \leq \bar{y}_{j'}\).

Similarly, the fact that the conditions (5) and (6) are not satisfied for the interval function \(F - v \cdot x\) mean that

\[ \exists i, j (i_0 \leq i < j \leq j_0 \& \bar{y}_i - v \cdot x_i \geq \bar{y}_j - v \cdot x_j) \]  \hfill (11)

and

\[ \exists i', j' (i_0 \leq i' < j' \leq j_0 \& \bar{y}_{i'} - v \cdot x_{i'} \leq \bar{y}_{j'} - v \cdot x_{j'}). \]  \hfill (12)

The inequality \(\bar{y}_j - v \cdot x_j \geq \bar{y}_j - v \cdot x_j\) can be described in the equivalent form \(v \cdot (x_j - x_i) \geq \bar{y}_j - \bar{y}_i\), i.e., since \(x_i < x_j\), in the form \(v \geq \bar{\Delta}_{ij}\). Thus, the existence of \(i\) and \(j\) as expressed by the formula (11) can be described as the existence of \(i\) and \(j\) for which \(v\) is larger than the corresponding value \(\bar{\Delta}_{ij}\), i.e., as

\[ v \geq \min_{i_0 \leq i < j \leq j_0} \bar{\Delta}_{ij}. \]

Similarly, the condition (12) is equivalent to

\[ v \leq \max_{i_0 \leq i < j \leq j_0} \Delta_{ij}. \]

The proposition is proven.

6 Towards a Faster Algorithm

Proposition 7 provides an explicit formula for computing \(F'([a, b])\) for each interval \([a, b]\). For each \([a, b]\), we need to compute \(O(n^2)\) values of \(\Delta_{ij}\) and \(\bar{\Delta}_{ij}\).

In problem like locating landscape features, we are not so much interested in knowing whether a given type of landscape exists in a given zone, but rather in locating all types of landscape. In other words, we would like to be able to find the values \(F'([a, b])\) for all possible intervals \([a, b]\). According to Proposition 7, it is sufficient to find all the values \(F'([x_{i_0}, x_{j_0}])\) for all \(i_0, j_0 = 1, \ldots, n\) for which \(i_0 < j_0\). There are \(n \cdot (n + 1)/2 = O(n^2)\) such values. If we use the formula from Proposition 7 – that takes \(O(n^2)\) computational steps – to compute each of these \(O(n^2)\) values, we will need an overall of \(O(n^2) \cdot O(n^2) = O(n^4)\) steps.

For large \(n – e.g., for n \approx 10^6\) – we need \(n^4 \approx 10^{24}\) computational steps; this is too long for even the fastest computers. Let us show that we can compute the interval derivative faster, actually in \(O(n^2)\) time. Since we must return \(O(n^2)\) results, we cannot do it in less than \(O(n^2)\) computational steps – so this algorithm is (asymptotically) optimal.

**Proposition 8.** There exists an algorithm that, given an interval function \(F = \{x_1, y_1, \ldots, x_n, y_n\}\), computes all possible values of the derivative \(F'([a, b])\) in \(O(n^2)\) computational steps.

**Proof.** At first, we compute \(O(n^2)\) values \(\Delta_{ij}\) and \(\bar{\Delta}_{ij}\) by using the formulas (9) and (10); this requires \(O(n^2)\) steps.

Let us now show how to compute all \(n^2\) values \(F_{i_0, j_0}\) in \(O(n^2)\) steps.
First, for each \( i \), we sequentially compute the “vertical” maxima \( \tau_{ij} \) corresponding to \( j = i + 1, i + 2, \ldots, n \) as follows:

\[
\tau_{i,j+1} = \Delta_{i,j+1} \quad \text{and} \quad \tau_{ij} = \max(\tau_{i,j-1}, \Delta_{ij}) \quad \text{for} \quad j > i + 1.
\]

For each \( i = 1, \ldots, n \), to compute all these values, we need \( n \) computational steps. Thus, to compute all such values \( \tau_{ij} \) for all \( i \) and \( j \), we need \( n \cdot n \) or \( O(n^2) \) computational steps.

Then, for every \( j_0 \), we sequentially compute the values \( F_{i_0,j_0} \) for \( i_0 = j_0 - 1, j_0 - 2, \ldots, 1 \) as follows: \( F_{j_0-1,j_0} = \Delta_{j_0-1,j_0} \) and \( F_{i_0,j_0} = \max(F_{i_0+1,j_0}, \tau_{i_0,j_0}) \) (it is easy to see that this formula is indeed correct). For each \( j_0 = 1, \ldots, n \), to compute all these values, we need \( n \) computational steps. Thus, to compute all such values \( F_{ij} \) for all \( i_0 \) and \( j_0 \), we need \( n \cdot n = O(n^2) \) computational steps.

Similarly, by using \( \Sigma_{ij} \) instead of \( \Delta_{ij} \) and \( \min \) instead of \( \max \), we can compute all \( n^2 \) values \( E_{0,j_0} \) in \( O(n^2) \) steps. The proposition is proven.

### 7 This Same Differential Formalism Also Serves an Alternative Definition of Zones

In some practical problems, a zone is defined not by an exact value of the derivative \( \nu \), but an interval \( \nu = [\overline{\nu}, \underline{\nu}] \) of possible values. In this case, it makes sense to say that an interval \( [a, b] \) contains a zone if for every function \( f \in F \), there is at least one point \( x \in [a, b] \) for which \( f'(x) \in \nu \). In other words, we say that the interval \( [a, b] \) contains a zone of a given type if \( f'([a, b]) \cap \nu \neq \emptyset \) for all functions \( f \in F \).

It turns out that the above notion of a derivative can help us detect such zones as well. Namely, the following statement is true:

**Proposition 9.** For every interval function \( F \) and for every two intervals \( [a, b] \) and \( \nu \), the following properties are equivalent to each other:

- for every function \( f \in F \), we have \( f'([a, b]) \cap \nu \neq \emptyset \);
- \( E_{0,j_0} > \overline{\nu} \) or \( F_{i_0,j_0} < \nu \).

Indeed, let us assume that there exists a function \( f \in F \) for which \( f'([a, b]) \cap \nu = \emptyset \). Since every function \( f \in F \) is continuously differentiable, its derivative \( f'(x) \) is a continuous function, hence the range \( f'([a, b]) \) is an interval. There are two possible situations when this interval range does not intersect with \( \nu \):

- either all the values from this range are \( > \overline{\nu} \),
- or all the values from this range are \( < \nu \).

In the first case, we have \( f'(x) > \overline{\nu} \) for all \( x \in [a, b] \). Therefore, for the function \( g(x) \) defined as \( f(x) - \nu \cdot x \), we get \( g'(x) > 0 \) for all \( x \), i.e., the function \( g(x) \) is strongly increasing. Since \( f \in F \), we have \( g \in G \) defined \( F - \nu \cdot x \). Due to Proposition 1, the existence of a strongly increasing function \( g \in G \) means that \( y_i - \nu \cdot x_i < y_j - \nu \cdot x_j \) for all \( i < j \). This inequality, in its turn, means that \( \Sigma_{ij} > \overline{\nu} \) for all \( i < j \). Thus, \( \overline{\nu} \) is smaller than the smallest of the values \( \Delta_{ij} \), i.e., smaller than \( E_{0,j_0} \).

Similarly, in the second case, we have \( f'(x) < \nu \) for all \( x \in [a, b] \), hence \( F_{i_0,j_0} < \nu \).

Vice versa, let \( E_{0,j_0} > \overline{\nu} \). By definition of \( E_{0,j_0} \) as the minimum, this means that \( \Sigma_{ij} > \overline{\nu} \) for all \( i, j \) for which \( i_0 \leq i < j \leq j_0 \). Substituting the definition of \( \Sigma_{ij} \), multiplying both sides of the inequality by a positive term \( x_j - x_i \) and moving terms to another side, we conclude that \( y_i - \nu \cdot x_i < y_j - \nu \cdot x_j \) for all \( i < j \). This inequality, in its turn, means that for the interval function \( G \) defined \( F - \nu \cdot x \), formula (2) holds and thus, due to Proposition 1, there exist a strongly monotonic function \( g \in G \) for which \( g'(x) > 0 \) for all \( x \). Then, for the function \( f(x) \) defined \( g(x) + \nu \cdot x \), we have \( f \in F \) and \( f'(x) = g'(x) + \nu \cdot x \), hence \( f'([a, b]) \cap \nu = \emptyset \).

Similarly, if \( F_{i_0,j_0} < \nu \), there exists a function \( f \in F \) for which \( f'([a, b]) \cap \nu = \emptyset \). The proposition is proven.

### 8 Open Problems

What if we take into consideration uncertainty in measuring \( x \)? In the above text, we took into consideration the uncertainty of measuring \( y \), but assumed that we know \( x \) exactly. In real life, there is
also some uncertainty in measuring $x$ as well. How can we take this uncertainty into consideration?

For the problem of finding local minima and maxima, this uncertainty was taken into consideration in [9]. It is desirable to extend this approach to finding the range of the derivatives.

**Parallelization.** In the above text, we described how to compute the derivative of an interval function in time $O(n^2)$, where $n$ is the number of observations, and showed that this algorithm is (asymptotically) optimal in the sense that no algorithm can compute this derivative faster.

For reasonable $n$, e.g., for $n \approx 10^3$, $n^2$ computational steps means a million steps; it is quite doable on modern computers. However, for large $n$, e.g., for $n \approx 10^6$, $n^2$ computational steps is $10^{12}$ steps, so on a modern Gigaherz machine, the corresponding computations will take $10^9$ sec – almost an hour.

How can we further speed up the corresponding computations? Our optimality result shows that we cannot achieve a drastic speed-up if we use sequential computers. Thus, the only way to speed up the corresponding computations is to use parallel computers.

For the problem of finding local minima and local maxima, parallel computers can indeed speed up the corresponding computations; see, e.g., [21]. An important question is therefore: How can speed up the computation of the corresponding derivative by using parallel computers?

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**References**


